

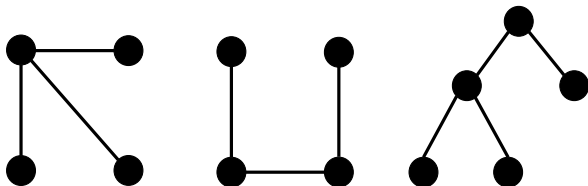
ET 1201  
Matematika Diskrit  
Prodi Teknik Telekomunikasi, Institut Teknologi Bandung  
2024

# Outline

- Tree
- Spanning Tree
- Huffman Code

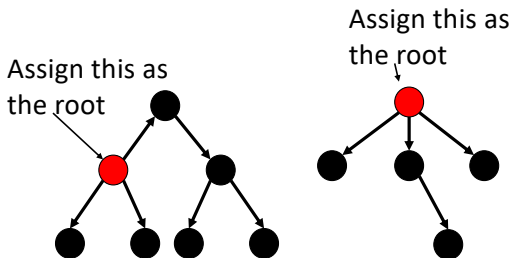
# What is a Tree ?

- ▶ An un-directed graph is a **tree** if and only if there is a unique simple path between any two of its vertices.
  - Only a **unique** simple path between two vertices.
  - No loops, no multiple edges.
  - A connected graph with no simple circuits.
- ▶ Examples of the tree.



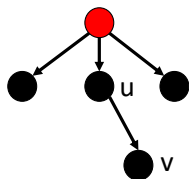
## Rooted Trees

- ▶ A **rooted tree** is:
  - One vertex has been designated as the root.
  - Every edge is directed away from the root.
- ▶ We usually put the root at the top, and point each edge downwards.

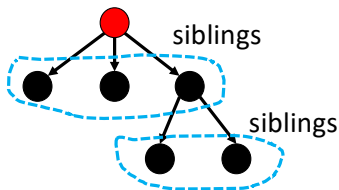


# Rooted Trees

- ▶ Each edge is from a **parent** to a **child**
  - The parent of a vertex is **unique**
- ▶ Vertices with the same parent are **siblings**



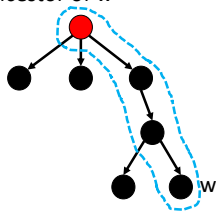
u is the parent of v  
v is a child of u



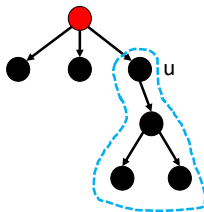
# Rooted Trees

- ▶ The **ancestors** of a vertex  $w$  include all the nodes in the path from the root to  $w$
- ▶ The **descendants** of a vertex  $u$  include all the nodes that have  $u$  as its ancestor.
- ▶ The **subtree** rooted at  $u$  includes all the descendants of  $u$ , and all edges that connect between them.

Each node is an ancestor of  $w$



Each node is an descendant of  $u$

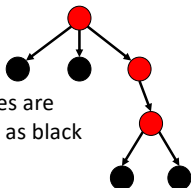


The whole part is the subtree rooted at  $u$

# Rooted Trees

- ▶ Vertices with no children are called **leaves**.
- ▶ Otherwise, they are called **internal nodes (vertices)**.
- ▶ A **m-ary** tree is every internal node has no more than  $m$  children.
- ▶ A **full m-ary tree** is every internal vertex has exactly  $m$  children.
  - A m-ary tree with  $m = 2$  is called **binary tree**.

All internal nodes  
are colored as red



All leaves are  
colored as black

The tree is ternary (3-ary),  
but not full

# Properties of Trees

► **Theorem:** A tree with  $n$  nodes has  $n - 1$  edges.

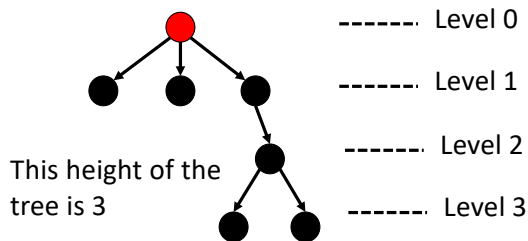
► **Proof:**

- Pick a vertex  $u$  in the tree and make  $u$  the root.
- Each edge links a parent and a child.
- The root has no parent, but every node has exactly one parent.
- Therefore, the number of edges =  $n - 1$ .



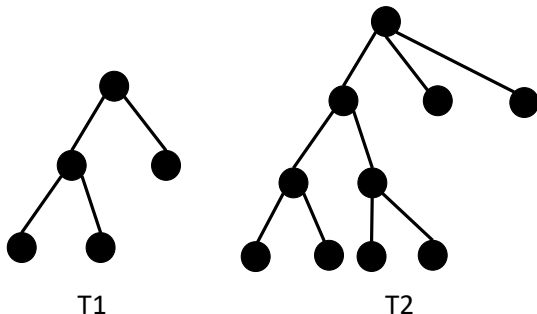
# Properties of Trees

- ▶ The **level** of vertex  $v$  in a rooted tree is the length of the path from root to  $v$
- ▶ The **height** of the tree is the maximum level of all the vertices.



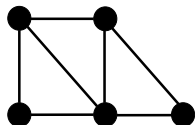
# Properties of Trees

- ▶ A rooted tree of height  $h$  is **balanced** if all the leaves are at level  $h$  or  $h - 1$
- ▶ Which of the following trees are balanced ?
  - T1 is balanced, because all its leaves are at level 2.
  - T2 is not balanced, because it has leaves at level 1 and level 3.

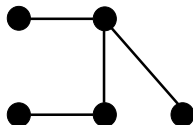


# What is a Spanning Tree ?

- ▶ A **spanning tree** of  $G$  is a subgraph of  $G$  that is a tree containing every vertex of  $G$ .



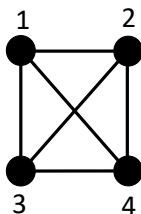
$G$



A spanning tree of  $G$

# Counting Spanning Trees

- ▶ Let  $G = (V, E)$ , and  $n$  is the number of vertices and  $m$  is the number of edges of  $G$ .  $K_n$  denotes a complete graph with  $n$  vertices.
- ▶ How many spanning trees are there in the complete graph  $K_n$ ?

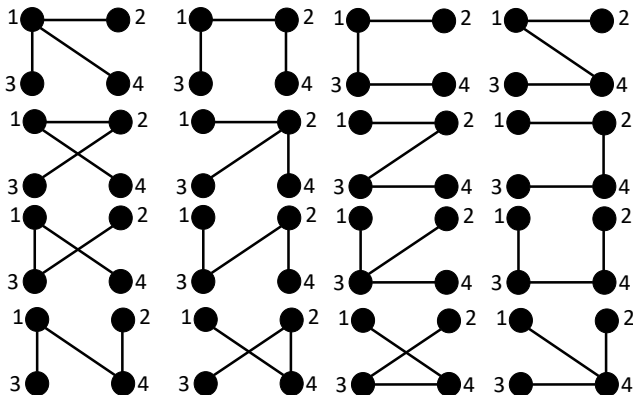


A four-vertex  
complete graph  $K_4$



# Counting Spanning Trees

- ▶ **16** spanning trees are in  $K_4$ .



# Text Encoding

- ▶ Each English character is represented in the same number of bits (8 bits) in ASCII.
  - ASCII uses **fixed-length** encoding
  - A text contains  $n$  characters, which take  $8n$  bits in total to store the text in ASCII.
  - Is it possible to find a coding scheme of these letters such that fewer bits are used ?

# Text Encoding

- ▶ In real-life English texts, characters do not appear with the same frequency.
- ▶ Using bit strings of different lengths to encode letters – **variable-length** encoding.
  - Frequent characters are encoded in fewer bits.
  - In-frequent characters are encoded in more bits.
- ▶ Then, we can reduce the total storage.

# Text Encoding

- ▶ Supposed a file contains 100K chars composed of A, B, C, D, E letters only.
  - A occurs 45K times, others 11K times each.
- ▶ Using fixed-length
  - Each character is encoded in 3 bits, total takes **300Kb**
- ▶ Using variable-length:
  - A  $\rightarrow$  0, B  $\rightarrow$  100, C  $\rightarrow$  101, D  $\rightarrow$  110, E  $\rightarrow$  111
  - $45K \times 1 + 44K = 177Kb$  (41% savings. )



# Prefix Code

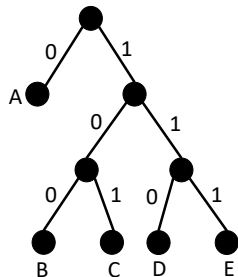
- ▶ If the encoding code book becomes:
  - $A \rightarrow 0, B \rightarrow 1, C \rightarrow 00, D \rightarrow 01, E \rightarrow 010$ .
  - Suppose the encoded text is: 0101
  - The original texts can have several possibilities such as
    - ABAB or ABD or DAB or DD or EB
- ▶ The problem comes from:
  - One codeword is a **prefix** of another.

# Prefix Code

- ▶ **Prefix code** encoding scheme is used to resolve that each codeword is a prefix of another.
- ▶ For a text encoded by a prefix code, we can easily decode it in the following way:
  - Scan from left to right to extract the first code
  - Recursively decode the remaining part.

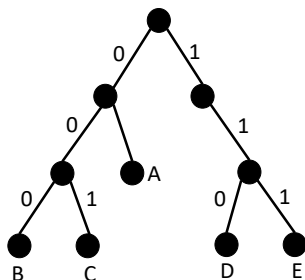
# Prefix Code Tree

- ▶ A **prefix code tree** is a rooted tree such that:
  - each edge is labeled by a bit
  - each leaf denoted by a character.
  - The codeword for the character is based on the labels on root-to-leaf path.
  - $A \rightarrow 0, B \rightarrow 100, C \rightarrow 101, D \rightarrow 110, E \rightarrow 111$



# Optimal Prefix Code

- ▶ **Problem:** Given the frequencies of each character, design the optimal prefix code whose encoded text requires the least storage.
- ▶ **Property 1:** In an optimal prefix code tree, each internal node must have two children.



- ▶ This is not an optimal prefix code tree

# Optimal Prefix Code

## ► **Property 2:**

- The leaves corresponding to the two least frequent characters are siblings.
- The leaves are farthest from the root.

## ► **Proof:** Consider an optimal prefix code tree.

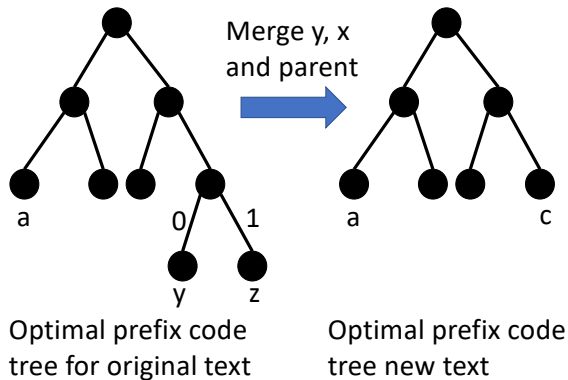
- Let  $y$  and  $z$  be the least frequent characters.
- Let  $x$  be a character whose leaf is the farthest from the root. Its sibling must be a leaf for some character  $x'$ .

# Optimal Prefix Code

- ▶ Let  $y$  and  $z$  be the two least frequent characters.
- ▶ Let  $T$  be an optimal tree such that  $y$  and  $z$  are sibling leaves and farthest from the root.
- ▶ When a new text shows: Replace each  $y$  and  $z$  by a common character  $c$  in the original text
- ▶ **Property 3:** We get an optimal prefix code tree for the new text if we merge  $y$ ,  $z$  and their parent into a leaf in  $T$ , and correspond this leaf to  $c$ .

# Optimal Prefix Code

- Graphically, the property 3 says:



# Optimal Prefix Code

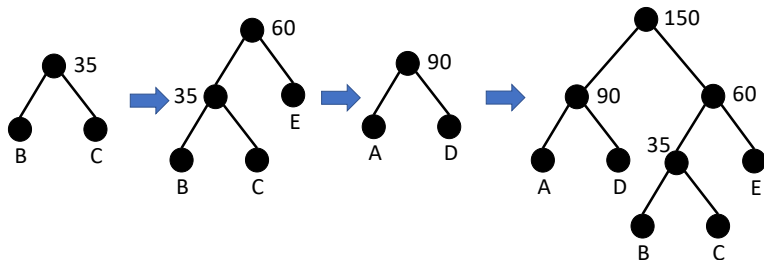
- ▶ Steps to obtain an optimal prefix code (David Huffman in 1952)
  - Find the least frequent characters  $x$  and  $y$ .
  - For two leaves for  $x$  and  $y$ , and join them with a common parent  $p$ .
  - Replace  $x$  and  $y$  by a common character  $c$ .
  - Recursively find the optimal prefix code tree for the new text (and replace the leaf for  $c$  with  $p$ ,  $x$ ,  $y$ ).



## Example

- Suppose the relative frequencies are as follows:

- A: 40, B: 20, C: 15, D: 50, E:25



# Referensi

- Lecture slides on DCP 1244 Discrete Mathematics, Tsung Tai Yeh, 2021

Available:

[https://people.cs.nycu.edu.tw/~ttyeh/course/2021\\_Spring/DCP1244/outline.html](https://people.cs.nycu.edu.tw/~ttyeh/course/2021_Spring/DCP1244/outline.html)