Assembly Programming II

CSE 351 Spring 2017

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Administrivia

- Lab 1 due Friday (4/14)
  - Remember, you have *late days* available if needed.
- Homework 2 due next Wednesday (4/19)

Three Basic Kinds of Instructions

1) Transfer data between memory and register
   - *Load* data from memory into register
     - `%reg` = Mem[address]
   - *Store* register data into memory
     - Mem[address] = `%reg`

Remember: Memory is indexed just like an array of bytes!

2) Perform arithmetic operation on register or memory data
   - `c = a + b;    z = x << y;    i = h & g;`

3) Control flow: what instruction to execute next
   - Unconditional jumps to/from procedures
   - Conditional branches

Operand types

- ***Immediate:*** Constant integer data
  - Examples: **$0x400** (hex), **$-533** (decimal)
  - Like C literal, but prefixed with **'$'**
  - Encoded with 1, 2, 4, or 8 bytes *depending on the instruction*
- ***Register:*** 1 of 16 integer registers
  - Examples: **%rax**, **%r13**
  - But **%rsp** reserved for special use
  - Others have special uses for particular instructions
- ***Memory:*** Consecutive bytes of memory at a computed address
  - Simplest example: **(%rax)** ← dereference: read data in %rax, treat as address, pull data from Mem starting at that address
  - Various other "address modes"

Registers
%rax
%rcx
%rdx
%rbx
%rsi
%rdi
%rsp (stack pointer)
%rbp
%rN (r8-r15)

Moving Data

AT&T syntax!

- General form: `mov_ source, destination`
  - Missing letter (_) specifies size of operands
  - Note that due to backwards-compatible support for 8086 programs (16-bit machines!), "word" means 16 bits = 2 bytes in x86 instruction names
  - Lots of these in typical code

- `movb src, dst`
  - Move 1-byte "**b**yte"
- `movw src, dst`
  - Move 2-byte "**w**ord"
- `movl src, dst`
  - Move 4-byte "**l**ong word"
- `movq src, dst`
  - Move 8-byte "**q**uad word"

movq Operand Combinations

	Source	Dest	Src, Dest	C Analog
movq	Imm	Reg	`movq $0x4, %rax`	`var_a = 0x4;`
		Mem	`movq $-147, (%rax)`	`*p_a = -147;`
	Reg	Reg	`movq %rax, %rdx`	`var_d = var_a;`
		Mem	`movq %rax, (%rdx)`	`*p_d = var_a;`
	Mem	Reg	`movq (%rax), %rdx`	`var_d = *p_a;`

- *Cannot do memory-memory transfer with a single instruction*
  - How would you do it?

x86-64 Introduction

- Arithmetic operations
- Memory addressing modes
  - `swap` example
- Address computation instruction (`lea`)

Some Arithmetic Operations

- Binary (two-operand) Instructions: Imm, Mem, Reg
  - **Maximum of one memory operand**
  - Beware argument order!
  - No distinction between signed and unsigned
    - Only arithmetic vs. logical shifts
  - How do you implement "`r3 = r1 + r2`"?

Format	Computation	
`addq src, dst`	*dst = dst + src*	(*dst += src*)
`subq src, dst`	*dst = dst – src*	
`imulq src, dst`	*dst = dst * src*	signed mult
`sarq src, dst`	*dst = dst >> src*	**A**rithmetic
`shrq src, dst`	*dst = dst >> src*	**L**ogical
`shlq src, dst`	*dst = dst << src*	(same as `salq`)
`xorq src, dst`	*dst = dst ^ src*	
`andq src, dst`	*dst = dst & src*	
`orq src, dst`	*dst = dst | src*	

operand size specifier

r3 = 0;
r3 += r1;
r3 += r2;

r3 = r1; ① movq r1, r3
r3 += r2; ② addq r2, r3

(src dst)

Some Arithmetic Operations

- Unary (one-operand) Instructions:

Format	Computation	
`incq` *`dst`*	*dst = dst + 1*	increment
`decq` *`dst`*	*dst = dst – 1*	decrement
`negq` *`dst`*	*dst = –dst*	negate
`notq` *`dst`*	*dst = ~dst*	bitwise complement

- See CSPP Section 3.5.5 for more instructions: `mulq`, `cqto`, `idivq`, `divq`

Arithmetic Example

Register	Use(s)
%rdi	1^{st} argument (x)
%rsi	2^{nd} argument (y)
%rax	return value

```
long simple_arith(long x, long y)
{
  long t1 = x + y;
  long t2 = t1 * 3;
  return t2;
}
```

```
y += x;
y *= 3;
long r = y;
return r;
```

```
simple_arith:
  addq    %rdi, %rsi
  imulq     $3, %rsi
  movq    %rsi, %rax
  ret
```

Example of Basic Addressing Modes

```
void swap(long *xp, long *yp)
{
  long t0 = *xp;
  long t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```

temp = *xp;
*xp = *yp;
*yp = temp;

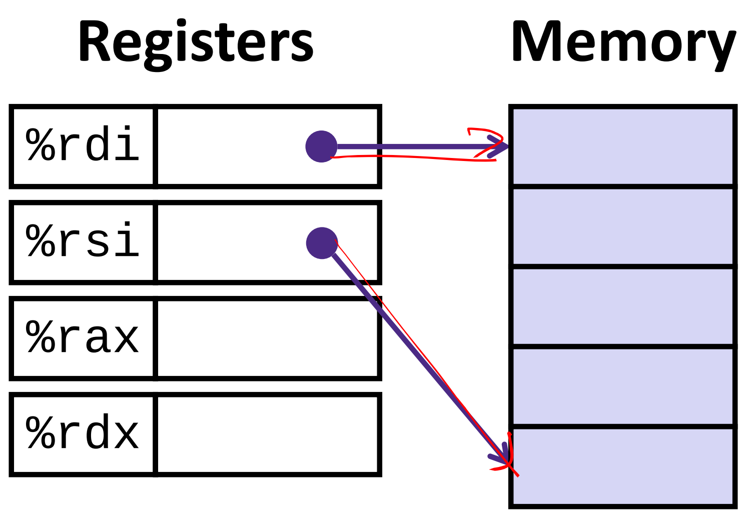
Mem → Reg
Reg → Mem

```
swap:
   movq  (%rdi), %rax
   movq  (%rsi), %rdx
   movq  %rdx, (%rdi)
   movq  %rax, (%rsi)
   ret
```

src dst
dereference!

Understanding swap()

```
void swap(long *xp, long *yp)
{
  long t0 = *xp;
  long t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```




```
swap:
   movq   (%rdi), %rax
   movq   (%rsi), %rdx
   movq   %rdx, (%rdi)
   movq   %rax, (%rsi)
   ret
```

Register		Variable
%rdi	⇔	xp
%rsi	⇔	yp
%rax	⇔	t0
%rdx	⇔	t1

Understanding swap()

Registers

%rdi	**0x120**
%rsi	**0x100**
%rax	
%rdx	

Memory	Word Address
123	0x120
	0x118
	0x110
	0x108
456	0x100

```
swap:
   movq   (%rdi), %rax   #  t0 = *xp
   movq   (%rsi), %rdx   #  t1 = *yp
   movq   %rdx, (%rdi)   # *xp =  t1
   movq   %rax, (%rsi)   # *yp =  t0
   ret
```

Understanding swap()

Registers

%rdi	0x120
%rsi	0x100
%rax	123
%rdx	

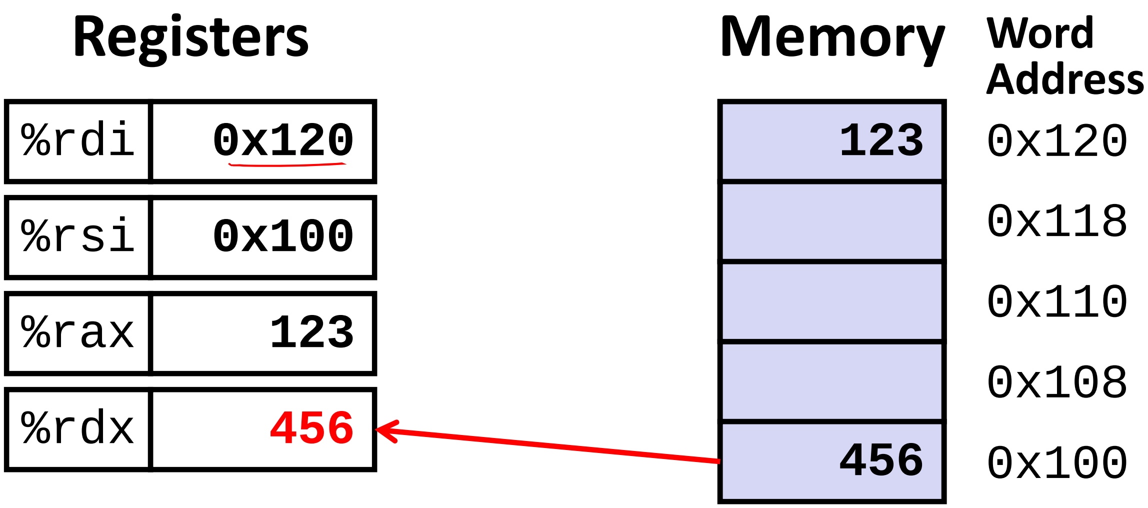
Memory

Memory	Word Address
123	0x120
	0x118
	0x110
	0x108
456	0x100

read %rdi; access that addr

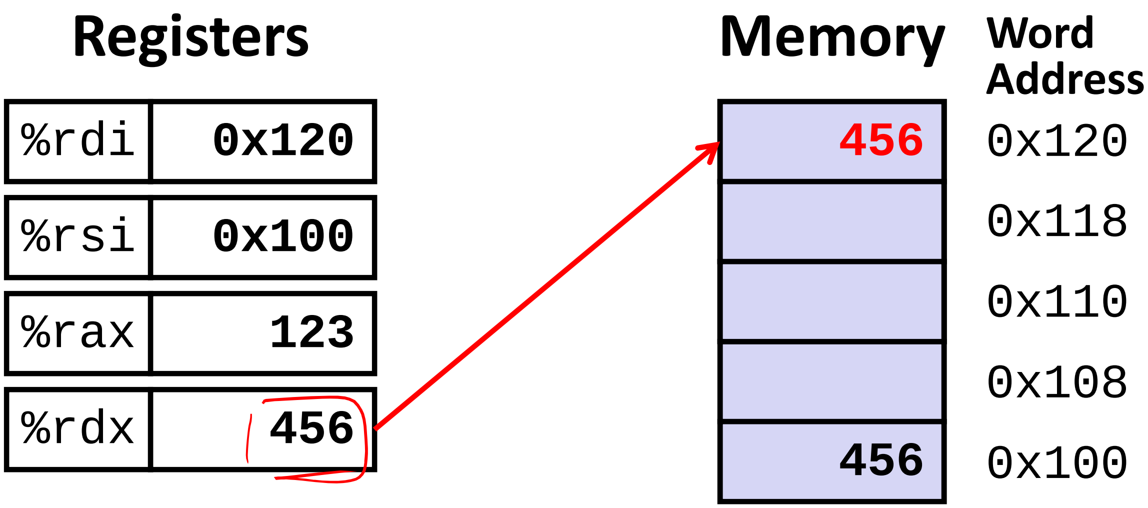
```
swap:
   movq   (%rdi), %rax   #  t0 = *xp
   movq   (%rsi), %rdx   #  t1 = *yp
   movq   %rdx, (%rdi)   # *xp =  t1
   movq   %rax, (%rsi)   # *yp =  t0
   ret
```

Understanding swap()



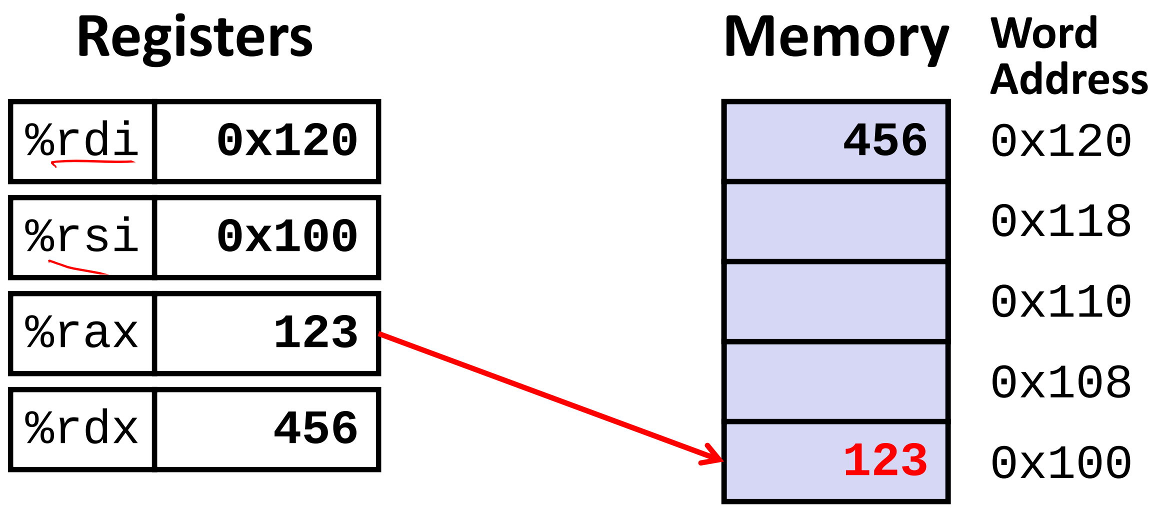

```
swap:
   movq   (%rdi), %rax   #  t0 = *xp
   movq   (%rsi), %rdx   #  t1 = *yp
   movq   %rdx, (%rdi)   # *xp =  t1
   movq   %rax, (%rsi)   # *yp =  t0
   ret
```

Understanding swap()




```
swap:
   movq   (%rdi), %rax   #  t0 = *xp
   movq   (%rsi), %rdx   #  t1 = *yp
   movq   %rdx, (%rdi)   # *xp =  t1
   movq   %rax, (%rsi)   # *yp =  t0
   ret
```

Understanding `swap()`




```
swap:
   movq   (%rdi), %rax   #  t0 = *xp
   movq   (%rsi), %rdx   #  t1 = *yp
   movq   %rdx, (%rdi)   # *xp =  t1
   movq   %rax, (%rsi)   # *yp =  t0
   ret
```

Memory Addressing Modes: Basic

treat as an array

- **Indirect:** (R) Mem[Reg[R]]

  take value in register

  - Data in register R specifies the memory address
  - Like pointer dereference in C
  - Example: **movq** (%rcx), %rax

- **Displacement:** D(R) Mem[Reg[R]+D]

  8

  - Data in register R specifies the *start* of some memory region
  - Constant displacement D specifies the offset from that address

    no space

  - Example: **movq** 8(%rbp), %rdx

Complete Memory Addressing Modes

ar[i] ⟷ *(ar+i) ⟶ Mem[ar + i* sizeof(data)]

- **General:**
  - D(Rb,Ri,S) Mem[Reg[Rb]+Reg[Ri]*S+D]
    - Rb: Base register (any register)
    - Ri: Index register (any register except %rsp)
    - S: Scale factor (1, 2, 4, 8) *– why these numbers?*
    - D: Constant displacement value (a.k.a. immediate)

- **Special cases** (see CSPP Figure 3.3 on p.181)
  - D(Rb,Ri) Mem[Reg[Rb]+Reg[Ri]+D] (S=1)
  - (Rb,Ri,S) Mem[Reg[Rb]+Reg[Ri]*S] (D=0)
  - (Rb,Ri) Mem[Reg[Rb]+Reg[Ri]] (S=1,D=0)
  - (,Ri,S) Mem[Reg[Ri]*S] (Rb=0,D=0)

↑ to differentiate Ri from Rb

Address Computation Examples

if not specified:
S = 1
D = 0
Reg[Rb] = 0
Reg[Ri] = 0

%rdx	0xf000
%rcx	0x0100

D(Rb,Ri,S) → Mem[Reg[Rb]+Reg[Ri]*S+D]

Expression	Address Computation	Address
0x8(%rdx)	Reg[Rb]+D = 0xf000 + 0x8	0xf008
(%rdx,%rcx)		0xf100
(%rdx,%rcx,4)	0xf000 + 4 * 0x100	0xf400
0x80(,%rdx,2)	0xf000 * 2 + 0x80	0x1E080

1111
1/1110

Address Computation Examples

%rdx	**0xf000**
%rcx	**0x0100**

D(Rb,Ri,S) → Mem[Reg[Rb]+Reg[Ri]*S+D]

Expression	Address Computation	Address
0x8(%rdx)	0xf000 + 0x8	0xf008
(%rdx,%rcx)	0xf000 + 0x100	0xf100
(%rdx,%rcx,4)	0xf000 + 0x100*4	0xf400
0x80(,%rdx,2)	0xf000*2 + 0x80	0x1e080

Address Computation Instruction

exception to the rule!

- `leaq src, dst`
  - "`lea`" stands for *load effective address*
  - `src` is address expression (any of the formats we've seen)
  - `dst` is a register
  - Sets `dst` to the *address* computed by the `src` expression (does not go to memory! – it just does math)
  - Example: `leaq (%rdx,%rcx,4), %rax` (addr)
- Uses:
  - Computing addresses without a memory reference
    - *e.g.* translation of `p = &x[i];`
  - Computing arithmetic expressions of the form `x+k*i+d`
    - Though k can only be 1, 2, 4, or 8

Example: `lea` vs. `mov`

Registers

%rax	
%rbx	
%rcx	**0x4**
%rdx	**0x100**
%rdi	
%rsi	

Memory

Memory	Word Address
0x400	0x120
0xF	0x118
0x8	0x110
0x10	0x108
0x1	0x100

```
leaq (%rdx,%rcx,4), %rax
movq (%rdx,%rcx,4), %rbx
leaq (%rdx), %rdi
movq (%rdx), %rsi
```

Example: `lea` vs. `mov` (solution)

Registers

%rax	0x110
%rbx	0x8
%rcx	0x4
%rdx	0x100
%rdi	0x100
%rsi	0x1

Memory	Word Address
0x400	0x120
0xF	0x118
0x8	0x110
0x10	0x108
0x1	0x100

```
leaq (%rdx,%rcx,4), %rax
movq (%rdx,%rcx,4), %rbx
leaq (%rdx), %rdi
movq (%rdx), %rsi
```

Arithmetic Example

Register	Use(s)
%rdi	1st argument (x)
%rsi	2nd argument (y)
%rdx	3rd argument (z)

```
long arith(long x, long y, long z)
{
  long t1 = x + y;
  long t2 = z + t1;
  long t3 = x + 4;
  long t4 = y * 48;
  long t5 = t3 + t4;
  long rval = t2 * t5;
  return rval;
}
```

```
arith:
  leaq    (%rdi,%rsi), %rax
  addq    %rdx, %rax
  leaq    (%rsi,%rsi,2), %rdx
  salq    $4, %rdx
  leaq    4(%rdi,%rdx), %rcx
  imulq   %rcx, %rax
  ret
```

- Interesting Instructions
  - `leaq`: "address" computation
  - `salq`: shift
  - `imulq`: multiplication
    - Only used once!

Arithmetic Example

Register	Use(s)
%rdi	x
%rsi	y
%rdx	z, t4
%rax	t1, t2, rval
%rcx	t5

```
long arith(long x, long y, long z)
{
  long t1 = x + y;
  long t2 = z + t1;
  long t3 = x + 4;
  long t4 = y * 48;
  long t5 = t3 + t4;
  long rval = t2 * t5;
  return rval;
}
```

```
arith:
  leaq    (%rdi,%rsi), %rax     # rax/t1   = x + y
  addq    %rdx, %rax            # rax/t2   = t1 + z
  leaq    (%rsi,%rsi,2), %rdx   # rdx      = 3 * y
  salq    $4, %rdx              # rdx/t4   = (3*y) * 16
  leaq    4(%rdi,%rdx), %rcx    # rcx/t5   = x + t4 + 4
  imulq   %rcx, %rax            # rax/rval = t5 * t2
  ret
```

Question

- Which of the following x86-64 instructions correctly calculates `%rax=9*%rdi`?

S ∈ {1,2,4,8}

A. `leaq (,%rdi,9), %rax`

B. `movq (,%rdi,9), %rax`

C. `leaq (%rdi,%rdi,8), %rax` → %rax = 9 * %rdi

D. `movq (%rdi,%rdi,8), %rax` → %rax = *(9 * %rdi)

x86 Control Flow

- Condition codes
- Conditional and unconditional branches
- Loops
- Switches

Control Flow

Register	Use(s)
%rdi	1st argument (x)
%rsi	2nd argument (y)
%rax	return value

```
long max(long x, long y)
{
  long max;
  if (x > y) {
    max = x;
  } else {
    max = y;
  }
  return max;
}
```

```
max:
  ???
  movq    %rdi, %rax
  ???
  ???
  movq    %rsi, %rax
  ???
  ret
```

Control Flow

Register	Use(s)
`%rdi`	1st argument (x)
`%rsi`	2nd argument (y)
`%rax`	return value

```
long max(long x, long y)
{
  long max;
  if (x > y) {
    max = x;
  } else {
    max = y;
  }
  return max;
}
```

Conditional jump

Unconditional jump

```
max:
  if x <= y then jump to else
  movq   %rdi, %rax
  jump to done
else:
  movq   %rsi, %rax
done:
  ret
```

Conditionals and Control Flow

- Conditional branch/*jump*
  - Jump to somewhere else if some *condition* is true, otherwise execute next instruction
- Unconditional branch/*jump*
  - *Always* jump when you get to this instruction

- Together, they can implement most control flow constructs in high-level languages:
  - **if** (*condition*) **then** {…} **else** {…}
  - **while** (*condition*) {…}
  - **do** {…} **while** (*condition*)
  - **for** (*initialization*; *condition*; *iterative*) {…}
  - **switch** {…}

Summary

- **Memory Addressing Modes:** The addresses used for accessing memory in `mov` (and other) instructions can be computed in several different ways
  - *Base register*, *index register*, *scale factor*, and *displacement* map well to pointer arithmetic operations
- `lea` is address calculation instruction
  - Does NOT actually go to memory
  - Used to compute addresses or some arithmetic expressions
- Control flow in x86 determined by status of Condition Codes